

Amendments to the Claims:

This listing of claims will replace all prior versions, and listings of claims in the application:

Listing of Claims:

1.-9. (canceled).

10. (currently amended) A method performed by a custodian computing system, having memory, to share a secret S among n secret owners such that any k of the n secret owners may reconstruct the secret, the method comprising the steps of:

choosing two large primes P and Q , such that PQ is greater than S ;

computing, at the custodian computing system, and storing in the custodian computer memory a product $N = PQ$;

computing and storing a product $M = (P-1)(Q-1)$;

choosing n random numbers e_1 through e_n that are relatively prime to N ;

choosing another random number e that is relatively prime to N ;

choosing n numbers d_1 through d_n such that $e_i d_i \bmod M$ equals one for $1 \leq i \leq n$;

choosing another number d such that $ed \bmod M$ is equal to one;

generating and storing a database of $\binom{n}{k}$ values, where each value is the product

of d and a unique k of the d_i numbers for $1 \leq i \leq n$, wherein each value is associated with a unique combination of k secret owners of the n secret owners;

deleting from the custodian computer memory P , Q , and M ;

computing S^e ;

distributing n secret owner pieces to each of the n secret owners, wherein each of the secret owner pieces includes S^e and one of the numbers e_1 through e_n ; and

deleting the secret S and e_1 through e_n , e , d_1 through d_n , and d ;

receiving k secret owner values from a unique combination of k secret owners;

determining a value c that is associated with the unique combination; and

determining the secret S using the value c and the k secret owner values.

11. (currently amended) A method as in claim 10, wherein receiving k secret owner values from the unique combination of k secret owners comprises:~~the method further comprising the steps of:~~

receiving a first of the n secret owner pieces from one of the n secret owners; and
computing and storing $S' = S^{e_1} \bmod N$, where f represents the one of the numbers e_1 through e_n contained in the first of the n secret owner pieces.

12. (currently amended) A method as in claim 11, wherein receiving k secret owner values from the unique combination of k secret owners comprises:~~the method further comprising the steps of:~~

receiving a second of the n secret owner pieces from another one of the n secret owners;

computing $S^q \bmod N$, where q represents the one of the numbers e_1 through e_n contained in the second of the n secret owner pieces; and replacing S' with $S^q \bmod N$.

13. (currently amended) A method as in claim 12, wherein receiving k secret owner values from the unique combination of k secret owners comprises~~further comprising the step of:~~

each time another of the secret owner pieces is received from another one of the n secret owners;

computing $S^q \bmod N$, where q represents the one of the numbers e_1 through e_n contained in the another of the n secret owner pieces; and replacing S' with $S^q \bmod N$.

14. (currently amended) A method as in claim 13, further comprising the steps of:

after k secret owner pieces have been received,

retrieving from the database ~~a~~the value c from among the $\binom{n}{k}$ values, wherein the value c corresponds to the k secret owner pieces of the unique combination of k secret owners that were received by the custodian;

computing $S^c \bmod N$; and
replacing S' with $S^c \bmod N$.

15. (currently amended) A method performed by a custodian computing system, having memory, to share a secret S among n secret owners such that any k of the n secret owners may reconstruct the secret, the method comprising the steps of:

choosing two large primes P and Q , such that PQ is greater than S ;
computing, at the custodian computing system, and storing in the custodian computer memory a product $N = PQ$;
computing and storing a product $M = (P-1)(Q-1)$;
choosing n random numbers e_1 through e_n that are relatively prime to N ;
choosing random numbers e and e' that are relatively prime to N ;
choosing n numbers d_1 through d_n such that $e_i d_i \bmod M$ equals one for $1 \leq i \leq n$;
choosing numbers d and d' such that $ed \bmod M$ is equal to one and such that $e'd' \bmod M$ is equal to one;

generating and storing a database of $\binom{n}{k}$ values, where each value is the product of d and a unique k of the d_i numbers for $1 \leq i \leq n$, wherein each value is associated with a unique combination of k secret owners of the n secret owners;

deleting from the custodian computer memory P , Q , and M ;
computing $S^{ee'}$;
distributing n secret owner pieces to each of the n secret owners, wherein each of the secret owner pieces includes $S^{ee'}$ and one of the numbers e_1 through e_n ; and
deleting the secret S and e_1 through e_n , e , d_1 through d_n , and d

receiving k secret owner values from a unique combination of k secret owners;
determining a value c that is associated with the unique combination; and
determining the secret S using the value c and the k secret owner value.

16. (currently amended) A method as in claim 15, wherein receiving k secret owner values from the unique combination of k secret owners comprises~~the method further comprising the steps of:~~

receiving a first of the n secret owner pieces from one of the n secret owners; and
computing and storing $S' = S^{e_1 f} \bmod N$, where f represents the one of the numbers e_1 through e_n contained in the first of the n secret owner pieces.

17. (currently amended) A method as in claim 16, wherein receiving k secret owner values from the unique combination of k secret owners comprises~~the method further comprising the steps of:~~

receiving a second of the n secret owner pieces from another one of the n secret owners;

computing $S^q \bmod N$, where q represents the one of the numbers e_1 through e_n contained in the second of the n secret owner pieces; and replacing S' with $S^q \bmod N$.

18. (currently amended) A method as in claim 17, wherein receiving k secret owner values from the unique combination of k secret owners comprises~~further comprising the step of:~~

each time another of the secret owner pieces is received from another one of the n secret owners;

computing $S^q \bmod N$, where q represents the one of the numbers e_1 through e_n contained in the another of the n secret owner pieces; and replacing S' with $S^q \bmod N$.

19. (currently amended) A method as in claim 18, further comprising the steps of:

after k secret owner pieces have been received,

retrieving from the database ~~a~~ the value c from among the $\binom{n}{k}$ values, wherein the value c corresponds to the k secret owner pieces from the unique combination of k secret owners that were received by the custodian;

computing $S^c \bmod N$;

replacing S' with $S^c \bmod N$;

computing $S^{d'} \bmod N$; and

replacing S' with $S^{d'} \bmod N$.

20. (currently amended) A method performed by a custodian computing system, having memory, to share a secret among n secret owners such that any k of the n secret owners may reconstruct the secret, the method comprising the steps of:

encrypting the secret so as to generate an encrypted secret;

deleting from the custodian computer memory the secret; and

performing a forward k out of n secret sharing algorithm on the encrypted secret

so as to generate n secret owner pieces;

storing a plurality of values associated with a plurality of unique combinations of k secret owners of the n secret owners;

distributing the n secret owner pieces to the n secret owners;

receiving k secret owner values from a unique combination of k secret owners;;

determining a value c that is associated with the unique combination;

performing a reverse k out of n secret sharing algorithm on the k secret owner pieces so as to recreate the encrypted secret using the value c ; and

decrypting the encrypted secret so as to recreate the secret.

21. - 24. (canceled).

25. (original) A method as in claim 20, wherein the step of performing a forward k out of n secret sharing algorithm includes the steps of:

dividing the encrypted secret into k pieces; and

performing n polynomial evaluations at n points of a degree- k polynomial using the k pieces of the encrypted secret as polynomial coefficients;

wherein each of the k secret owner pieces includes a result of one of the n polynomial evaluations and a corresponding one of the n points.

26. (currently amended) A method as in claim 25, ~~further comprising the steps of:~~

~~distributing the n secret owner pieces to the n secret owners;~~
~~receiving k secret owner pieces from k secret owners; and~~
~~performing a reverse k out of n secret sharing algorithm on the k secret owner pieces so as to recreate the encrypted secret;~~ wherein the step of performing a reverse k out of n secret sharing algorithm includes the steps of generating a system of k linear equations and solving the system of k linear equations for the k pieces of the encrypted secret.

27. (currently amended) A method as in claim 26, further comprising the step of:
assembling the k pieces of the encrypted secret so as to recreate the encrypted
secret; and

~~decrypting the encrypted secret so as to recreate the secret.~~

28.-29. (canceled).

30. (currently amended) A computer readable storage medium having embodied thereon computer readable program code suitable for programming a computer to perform a method performed by a custodian to share a secret S among n secret owners such that any k of the n secret owners may reconstruct the secret, the method comprising the steps of:

choosing two large primes P and Q , such that PQ is greater than S ;

computing and storing a product $N = PQ$;

computing and storing a product $M = (P-1)(Q-1)$;

choosing n random numbers e_1 through e_n that are relatively prime to N ;

choosing another random number e that is relatively prime to N ;

choosing n numbers d_1 through d_n such that $e_i d_i \bmod M$ equals one for $1 \leq i \leq n$;

choosing another number d such that $ed \bmod M$ is equal to one;

generating and storing a database of $\binom{n}{k}$ values, where each value is the product

of d and a unique k of the d_i numbers for $1 \leq i \leq n$, wherein each value is associated with a unique combination of k secret owners of the n secret owners;

deleting P , Q , and M ;

computing S^e ;

distributing n secret owner pieces to each of the n secret owners, wherein each of the secret owner pieces includes S^e and one of the numbers e_1 through e_n ; and

deleting the secret S and e_1 through e_n , e , d_1 through d_n , and d ;

receiving k secret owner values from a unique combination of k secret owners;

determining a value c that is associated with the unique combination; and

determining the secret S using the value c and the k secret owner values.

31. (currently amended) A computer readable storage medium having embodied thereon computer readable program code suitable for programming a computer to perform a method performed by a custodian to share a secret S among n secret owners such that any k of the n secret owners may reconstruct the secret, the method comprising the steps of:

choosing two large primes P and Q , such that PQ is greater than S ;

computing and storing a product $N = PQ$;

computing and storing a product $M = (P-1)(Q-1)$;

choosing n random numbers e_1 through e_n that are relatively prime to N ;

choosing random numbers e and e' that are relatively prime to N ;

choosing n numbers d_1 through d_n such that $e_i d_i \bmod M$ equals one for $1 \leq i \leq n$;

choosing numbers d and d' such that $ed \bmod M$ is equal to one and such that $e'd' \bmod M$ is equal to one;

generating and storing a database of $\binom{n}{k}$ values, where each value is the product of d and a unique k of the d_i numbers for $1 \leq i \leq n$, wherein each value is associated with a unique combination of k secret owners of the n secret owners;

deleting P , Q , and M ;

computing $S^{ee'}$;

distributing n secret owner pieces to each of the n secret owners, wherein each of the secret owner pieces includes $S^{ee'}$ and one of the numbers e_1 through e_n ; and

deleting the secret S and e_1 through e_n , d_1 through d_n , and d_2 ;

receiving k secret owner values from a unique combination of k secret owners;

determining a value c that is associated with the unique combination; and

determining the secret S using the value c and the k secret owner values.

32. (currently amended) A computer readable storage medium having embodied thereon computer readable program code suitable for programming a computer to perform a method performed by a custodian to share a secret among n secret owners such that any k of the n secret owners may reconstruct the secret, the method comprising the steps of:

encrypting the secret so as to generate an encrypted secret;

deleting the secret; and

performing a forward k out of n secret sharing algorithm on the encrypted secret so as to generate n secret owner pieces;

storing a plurality of values associated with a plurality of unique combinations of k secret owners of the n secret owners;

distributing the n secret owner pieces to the n secret owners;

receiving k secret owner values from a unique combination of k secret owners;

determining a value c that is associated with the unique combination;

performing a reverse k out of n secret sharing algorithm on the k secret owner pieces so as to recreate the encrypted secret using the value c ; and

decrypting the encrypted secret so as to recreate the secret.

33.-34. (canceled).

35. (currently amended) A computer comprising a processor and a computer readable storage medium coupled to the processor having embodied thereon processor readable program code suitable for programming a computer to perform a method performed by a custodian to share a secret S among n secret owners such that any k of the n secret owners may reconstruct the secret, the method comprising the steps of:

choosing two large primes P and Q , such that PQ is greater than S ;

computing and storing a product $N = PQ$;

computing and storing a product $M = (P-1)(Q-1)$;

choosing n random numbers e_1 through e_n that are relatively prime to N ;

choosing another random number e that is relatively prime to N ;

choosing n numbers d_1 through d_n such that $e_i d_i \bmod M$ equals one for $1 \leq i \leq n$;

choosing another number d such that $ed \bmod M$ is equal to one;

generating and storing a database of $\binom{n}{k}$ values, where each value is the product

of d and a unique k of the d_i numbers for $1 \leq i \leq n$, wherein each value is associated with a unique combination of k secret owners of the n secret owners;

deleting P , Q , and M ;

computing S^e ;

distributing n secret owner pieces to each of the n secret owners, wherein each of the secret owner pieces includes S^e and one of the numbers e_1 through e_n ; and

deleting the secret S and e_1 through e_n , e , d_1 through d_n , and d ;

receiving k secret owner values from a unique combination of k secret owners;

determining a value c that is associated with the unique combination; and

determining the secret S using the value c and the k secret owner values.

36. (currently amended) A computer comprising a processor and a computer readable storage medium coupled to the processor having embodied thereon processor readable program code suitable for programming the computer to perform a method performed by a custodian to share a secret S among n secret owners such that any k of the n secret owners may reconstruct the secret, the method comprising the steps of:

- choosing two large primes P and Q , such that PQ is greater than S ;
- computing and storing a product $N = PQ$;
- computing and storing a product $M = (P-1)(Q-1)$;
- choosing n random numbers e_1 through e_n that are relatively prime to N ;
- choosing random numbers e and e' that are relatively prime to N ;
- choosing n numbers d_1 through d_n such that $e_i d_i \bmod M$ equals one for $1 \leq i \leq n$;
- choosing numbers d and d' such that $ed \bmod M$ is equal to one and such that $e'd' \bmod M$ is equal to one;

generating and storing a database of $\binom{n}{k}$ values, where each value is the product

of d and a unique k of the d_i numbers for $1 \leq i \leq n$, wherein each value is associated with a unique combination of k secret owners of the n secret owners;

deleting P , Q , and M ;

computing $S^{ee'}$;

distributing n secret owner pieces to each of the n secret owners, wherein each of the secret owner pieces includes $S^{ee'}$ and one of the numbers e_1 through e_n ; and

deleting the secret S and e_1 through e_n , e , d_1 through d_n , and d ;

receiving k secret owner values from a unique combination of k secret owners;

determining a value c that is associated with the unique combination; and

determining the secret S using the value c and the k secret owner values.

37. (currently amended) A computer comprising a processor and a computer readable storage medium coupled to the processor having embodied thereon processor readable

program code suitable for programming the computer to perform a method performed by a custodian to share a secret among n secret owner such that any k of the n secret owners may reconstruct the secret, the method comprising the steps of:

encrypting the secret so as to generate an encrypted secret;

deleting the secret; and

performing a forward k out of n secret sharing algorithm on the encrypted secret

so as to generate n secret owner pieces;

storing a plurality of values associated with a plurality of unique combinations of k secret owners of the n secret owners;

distributing the n secret owner pieces to the n secret owners;

receiving k secret owner values from a unique combination of k secret owners;

determining a value c that is associated with the unique combination;

performing a reverse k out of n secret sharing algorithm on the k secret owner pieces so as to recreate the encrypted secret using the value c ; and

decrypting the encrypted secret so as to recreate the secret.